## INFORMED SEARCH

Lara J. Martin (she/they) TA: Aydin Ayanzadeh (he)

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CMSC 671

## ADDITION TO PAPER SUMMARY

(no points but required at the end of the summary) the names of anyone you talked to about the paper and/or how you used LLMs (if you did). Please see the Generative AI Policy for more information on how to cite LLM use. If you did not receive any help, please state that instead.

## UNINFORMED MODELS COMPARED

## Breadth-First Search Depth-First Search Iterative Deepening

- Complete
- Optimal
- $O\left(b^{d}\right)$ time complexity
- $O\left(b^{d}\right)$ space complexity
- Not complete
- Not optimal
- $O\left(b^{m}\right)$ time complexity
- O(bm) space complexity
- Complete
- Optimal
- $O\left(b^{d}\right)$ time complexity
- O(bd) space complexity


# What makes a search algorithm "informed"? 

## INFORMED SEARCH INTRO

- down/right states are closer to the goal than the up/left states
- Define a function to represent approximate distance to the goal state: a heuristic function $h(n)$
- Example: Euclidean (straightline) distance
- Use $h(n)$ to select nodes for expansion



## RECAP

## Best-First Search uses what type of data structure for its frontier?

Priority queue

## A* SEARCH

## A* SEARCH

We need an optimal informed search algorithm

- Account for previous path length
- Uniform-Cost Search $g(n)$
- Account for approximate remaining path length
- Best-First Search $h(n)$
- Can we combine them?


## A* SEARCH

function TREE-SEARCH ( problem) returns a solution, or failure initialize the frontier using the initial state of problem

## loop do

if the frontier is empty then return failure
choose a leaf node and remove it from the frontier
if the node contains a goal state then return the corresponding solution expand the chosen node, adding the resulting nodes to the frontier

Store frontier as a priority queue, prioritized by $f(n)=g(n)+h(n)$

- $f(n)$ approximates the full solution cost passing through a node $n$
- $g(n)$ is the path cost and $h(n)$ is a heuristic approximating the cost-to-go


## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[Sibiu (0+253)]
Current node:
none


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[Sibiu (253)]
Current node:
none


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[]

Current node:
Sibiu (253)


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[R.V. (80+193), F. (99+176)]

Current node:
Sibiu (253)


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[R.V. (273), F. (275)]

Current node:
Sibiu (253)


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[F. (275)]
Current node: R.V. (273)


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[F. (275), P. (80+97+100)]
Current node:
R.V. (273)


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[F. (275), P. (277)]
Current node:
R.V. (273)


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[P. (277)]
Current node:
F. (275)


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[P. (277), B. $(99+211+0)$ ]
Current node:
F. (275)


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[P. (277), B.(310)]
Current node:
F. (275)


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[B.(310)]
Current node:
P. (277)


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[B. (80+97+101+0), B.(310)]
Current node:
P. (277)


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[B.(278), B.(310)]
Current node:
P. (277)


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest

## A* SEARCH EXAMPLE

Problem: find the shortest (in terms of distance) path from Sibiu to Bucharest
Frontier:
[B.(310)]
Current node:
B. (278),


| State | $\mathbf{h}$ (State) |
| :--- | :--- |
| Bucharest | 0 |
| Fagaras | 176 |
| Pitesti | 100 |
| Rimnicu <br> Vilcea | 193 |
| Sibiu | 253 |

Bucharest
Return:
[Sibiu, Rimnicu Vilcea, Pitesti, Bucharest]

## A* SEARCH PROPERTIES

## A* Search

- Complete
- Optimal ${ }^{1}$
- Exponential time complexity (decreased by good heuristics)
- Exponential space complexity (decreased by good heuristics)
- Optimally efficient ${ }^{1}: \mathrm{A}^{*}$ guaranteed to find optimal solution faster than any other informed search algorithm for a given heuristic $h(n)$
${ }^{1}$ For admissible heuristics (consistent heuristics for graph search)


## ADMISSIBLE HEURISTICS

## ADMISSIBLE HEURISTICS



$$
f^{*}(n)=g^{*}(n)+h^{*}(n)
$$

## Admissible heuristic: $\boldsymbol{h}(\boldsymbol{n}) \leq \boldsymbol{h}^{*}(\boldsymbol{n})$

$$
f(n)=g(n)+h(n) \leq C_{\text {actual }}^{*}
$$

A heuristic is consistent if it meets the triangle inequality:


## A* SEARCH PROPERTIES

If $h(n)$ is admissible: $f(n)=g(n)+h(n) \leq C^{*}$

- $\mathrm{A}^{*}$ is optimal because it expands all nodes where $f(n)<C^{*}$, and some nodes where $f(n)=C^{*}$
- $\mathrm{A}^{*}$ is efficient because it prunes (ignore parts of search space) where $f(n)>C^{*}$

Some heuristics are better than others:

- What happens if $h(n)=0$ ?
- Equivalent to UCS (still optimal, but uninformed)
- What happens if $h(n)=h^{*}(n)$, i.e. it estimates the exact cost to the goal?
- Every step will follow the optimal path


## EXAMPLES OF (SPATIAL) HEURISTICS

- Manhattan (city block) distance: sum of horizontal and vertical distance, minimum distance on a 4 -connected grid

$$
h(n)=\left|g_{x}-n_{x}\right|+\left|g_{y}-n_{y}\right|
$$

- Euclidean (straight-line) distance: shortest distance from start to goal

$$
h(n)=\sqrt{\left(g_{x}-n_{x}\right)^{2}+\left(g_{y}-n_{y}\right)^{2}}
$$

- Octile distance: minimum distance on a grid with diagonal moves (8-connected)


## A* SEARCH IN PRACTICE

Experiment URL (try it for yourself!):
https://qiao.github.io/PathFinding.js/visual/

Experiment summary:

- $\mathrm{A}^{*}$ succeeds where Best-First Search fails and UCS is inefficient
- Only holds for admissible heuristics
- Inadmissible heuristics may produce sub-optimal paths
- In the limit as $h(n)$ increases, $\mathrm{A}^{*}$ behaves like greedy search - $h(n)$ dominates $g(n)$, so $f(n) \approx h(n)$



## IS THIS HEURISTIC ADMISSIBLE? ROUTE FINDING

State: agent location (which path intersection are you at)
Actions: follow path $x$ to the next intersection Goal: find the shortest path to the end location Heuristic: Euclidean distance

Answer: Yes, because the shortest distance between two points is a line


## IS THIS HEURISTIC ADMISSIBLE? ROUTE FINDING

State: agent location (which path intersection are you at)
Actions: follow path $x$ to the next intersection Goal: find the shortest path to the end location Heuristic: Manhattan distance

Answer: No (e.g. it will overestimate diagonal paths)


## IS THIS HEURISTIC ADMISSIBLE? ROUTE FINDING



Heuristic: Manhattan distance (maybe yes if you're in Manhattan)

## IS THIS HEURISTIC ADMISSIBLE? 8-PUZZLE

State: configuration of tiles
Actions: slide any single tile to the open position
Goal: reach solved state in the fewest moves
Heuristic: \# of tiles out of solved position
Answer: Yes (we can only move one tile at a time, and any out of place tile will need to be moved at least once)

| $8^{X}$ |  | $6^{X}$ |
| :--- | :--- | :--- |
| $5^{X}$ | $\mathbf{4}^{\prime}$ | $\mathbf{7}^{X}$ |
| $2^{X}$ | $3^{X}$ | $\mathbf{1}^{x}$ |


|  | 1 |  |  | 2 |
| :---: | :---: | :---: | :---: | :---: |
| 3 | 4 | 5 |  |  |
| 6 | 7 | 8 |  |  |

## IS THIS HEURISTIC ADMISSIBLE? 8-PUZZLE

State: configuration of tiles
Actions: slide any single tile to the open position
Goal: reach solved state in the fewest moves
Heuristic: Manhattan distance for each tile to solved position
Answer: Yes (each tile can only move one space at a time, on a 4 -connected grid)

| 8 |  | 6 |
| :--- | :--- | :--- |
| 5 | 4 | 7 |
| 2 | 3 | 1 |



## IS THIS HEURISTIC ADMISSIBLE? RUBIK'S CUBE

State: configuration of blocks on the cube
Actions: rotate each face clockwise/counterclockwise 90 degrees
Goal: reach solved state in the fewest moves
Heuristic: \# of blocks out of solved position
Answer: No ( 1 single move can move 8 blocks at a time, so this can over estimate)


## WHICH HEURISTIC IS BETTER? 8-PUZZLE

Heuristic 1: \# of tiles out of solved position
Heuristic 2: Manhattan distance for each tile to solved position

Answer: Heuristic 2, because it dominates Heuristic 1 (see later slides)

| 8 |  | 6 |
| :--- | :--- | :--- |
| 5 | 4 | 7 |
| 2 | 3 | 1 |


|  | 1 |  |  | 2 |
| :---: | :---: | :---: | :---: | :---: |
| 3 | 4 | 5 |  |  |
| 6 | 7 | 8 |  |  |

## EVALUATING HEURISTICS

$A^{*}$ performance depends on how good your heuristic is.
How do we know how good a heuristic is?

- Informedness: a more informed heuristic will explore a smaller portion of the state space

$$
\operatorname{informedness}(h)=\frac{\# \text { states in search space }}{\text { avg \# of states explored with } h}
$$

- Dominance: a heuristic $h_{1}$ dominates $h_{2}$ if

$$
h_{2}\left(s_{i}\right) \leq h_{1}\left(s_{i}\right) \leq h^{*}\left(s_{i}\right) \text { for all states } s_{i}
$$

## OPTIMAL SEARCH ALGORITHM CHEATSHEET

| Search Algorithm | Frontier <br> Implementation | Notes |
| :---: | :---: | :---: |
| Breadth-First Search <br> $($ BFS $)$ | Queue |  |
| Depth-First Search <br> (DFS) | Stack | Optimal with <br> iterative deepening <br> Optimal with |
| Depth-Limited Search <br> Uniform-Cost Search <br> (UCS) | Priority Queue | Path cost $g(n)$ <br> known solution depth |
| A* Search | Priority Queue | Estimated solution cost <br> $f(n)=g(n)+h(n)$ |

## FOR NEXT CLASS

- Finish reading Chapter 3.5-3.7
- If you have Module $1 /$ Search for your paper presentation
- Summaries are due tomorrow!
- Presentation on Thursday Sept 21
- You should be ready to do all of HW 1 now

